

THE PROBLEM OF THE  $n$  QUEENS

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1. Introduction. The problem of the  $n$  queens is to determine in how many ways it is possible to place  $n$  queens on an  $n \times n$  chessboard so that no queen is attacking any other.

The original problem for  $n=8$  was first proposed by Max Bezzel in the Berlin Schachzeitung, in 1848, and the twelve solutions, by Franz Nauck, were published in 1850 in the Leipzig Illustrierte Zeitung. Since then a considerable literature has grown up around the problem and solutions for  $n \leq 12$  have been given in various places (Ahrens [1, vol 1, ch 9], Kraitchik [4, ch 10], Madachy [6, ch 2], Rouse Ball [7, ch 6]).

More recently interest has been shown in the problem as an exercise in computer programming and Dijkstra in [3, ch 1] uses it to introduce recursion as a programming technique. By developing the ideas of Dijkstra to search for a more restricted set of solutions it has been possible to extend the known results beyond those given by Kraitchik in [4, ch 10].

These results are given in Table I.

2. Definitions. A queen is defined as a pair of integers  $(a, b)$ , where  $1 \leq a \leq n$  and  $1 \leq b \leq n$ .

A solution,  $\omega$ , is a set of  $n$  queens.

$\{q_i | q_i = (i, b_i), i = 1, 2, \dots, n\}$ , where for all  $i$  and  $j$  satisfying  $1 \leq j < i \leq n$ , then  $|b_i - b_j| \neq i - j$ .

The norm of  $\omega$  is the base  $n$  number  $b_1 b_2 \dots b_n$ .

3. Generating the solutions by computer. Dijkstra in [3, ch 1] discusses a method for generating all solutions using a recursive technique. Solutions are generated more quickly, however, if the solutions are collected into classes, the members of which are equivalent under a group of transformations. The choice of the group requires that the equivalence class of any solution can be determined easily together with the number of solutions belonging to that class. A class can be identified by a characteristic solution which is defined as that solution, in the class, with the least norm. It is then sufficient to program a search for characteristic solutions only.

The number of solutions is denoted by  $T_n$  and the number

of characteristic solutions by  $C_n$

To determine  $T_n$  we choose the dihedral group of the square as the group of transformations. This group consists of the 8 transformations  $\{I, R, R^2, R^3, RS, R^2S, R^3S\}$ , where  $I$  denotes the identity transformation,  $R$  a rotation of the board in its own plane through  $\pi/2$  and  $S$  a reflection of the board. This group enables solutions to be classified as equivalent under simple rotations and reflections of the board (Alexandroff [2, p50], Kraitchik [4, ch 10], Ledermann [5, pp48-55]).

The equivalence class containing a solution  $\omega$  consists of:

- i) 8 members if  $\omega \neq R^2 \cdot \omega$  {  $\omega$  is unsymmetric }
- ii) 4 members if  $\omega = R^2 \cdot \omega$  but  $\omega \neq R \cdot \omega$  {  $\omega$  is symmetric }
- iii) 2 members if  $\omega = R \cdot \omega$  {  $\omega$  is doubly-symmetric }

The number of unsymmetric characteristic solutions is denoted by  $U_n$ , the number of symmetric characteristic solutions by  $S_n$  and the number of doubly-symmetric characteristic solutions by  $D_n$ .

Thus:

$$C_n = U_n + S_n + D_n$$

$$T_n = 8U_n + 4S_n + 2D_n$$

$n$	$P_n$	$D_n$	$S_n$	$U_n$	$C_n$	$T_n$
4	1	1	0	0	1	2
5	1	1	0	1	2	10
6	0	0	1	0	1	4
7	0	0	2	4	6	40
8	0	0	1	11	12	92
9	0	0	4	42	46	352
10	0	0	3	89	92	724
11	0	0	12	329	341	2680
12	1	4	18	1765	1787	14200
13	1	4	32	9197	9233	73712
14	0	0	105	45647	45752	395596
15	0	0	310			
16	4	32	734			
17	8	64	2006			
18	0	0	4526			
19	0	0				
20	15	240				
21	22	352				
22	0	0				
23	0	0				
24	52	1664				

25

51

1632

26	0	0
27	0	0
28	257	<del>21888</del> 16448
29	342	21888
30	0	0
31	0	0
32	1589	203392
	2609	333952
34	0	0
35	0	0
36	11417	2922752
37	16896	4325376
38	0	0
39	0	0
40	75375	38592000
41	99114	50746368