# A new approach to the middle levels via a Catalan-number system of numeration

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#### Abstract

Let n = 2k + 1 with  $0 < k \in \mathbb{Z}$  and let  $M_k$  be the subgraph induced by the kand (k + 1)-levels of the *n*-cube graph. A system of numeration according to which the nonnegative integers are written as restricted growth strings, so called by Arndt and Ruskey, has specifically the k-th Catalan number  $C_k = \binom{n}{k}/n$  expressed as  $10^k$ . This permits a linear ordering of the vertex set of a quotient graph  $R_k$  of  $M_k$  under a dihedral-group action. The Shields-Savage lifting of a Hamilton path in  $R_k$  between certain two distinguished vertices of  $R_k$  to a Hamilton cycle placed in a dihedrally symmetric fashion in  $M_k$  is presented in terms of the Kierstead-Trotter lexical matchings of  $M_k$ , which highlights the relevance of the said system in pursuing the existence, enumeration and sorting (according both to that system and to the lexical matchings) of such Hamilton cycles in  $M_k$ . With these tools in hand, we ask whether there are as many such cycles as the  $2^{2^{\Omega(n)}}$  ones obtained by Mütze in proving the existence of Hamilton cycles (not necessarily placed in dihedrally symmetric fashion) in  $M_k$ .

#### 1 Introduction

Let n = 2k + 1 with  $0 < k \in \mathbb{Z}$ . Assume that the dihedral group  $D_{2n}$  acts on a graph G, and in that case let H be a  $D_{2n}$ -invariant subgraph of G ([11], pg. 20). If this subgraph H is a Hamilton cycle of G, then we say that H is a Hamilton cycle placed in a dihedrally symmetric fashion in G.

In this work, each nonnegative integer is expressed as a restricted growth string (or RGS)  $\alpha$ , so called by Arndt [2] and Ruskey [17], (defined in Section 2, below). The resulting RGS s are related to the Catalan numbers  $C_k = \frac{1}{n} {n \choose k}$ , (A000108 in, [21]), in that  $C_k$  becomes the RGS  $\alpha = 10^k$ , as observed by Arndt in [2] pg. 325. A system of numeration, (A239903 in [21]), composed by the RGS s  $\alpha$  is presented in Section 2 and is taken to a dihedrally symmetric version of Hável's [12] or Buck-Wiedemann's [4] conjecture, proved by Mütze [16], on the existence of Hamilton cycles (not necessarily placed in a dihedrally symmetric fashion) in the middle-levels graphs. These graphs, denoted  $M_k$ , ( $0 < k \in \mathbb{Z}$ ), are treated

from Section 4 on, after associating the RGS s  $\alpha$  to corresponding *n*-tuples  $F(\alpha)$  representing the vertices of  $M_k$  by means of a "descending castling" operation in Theorem 2 of Section 3. In turn, the *n*-tuples  $F(\alpha)$  become, in Section 9, the vertices of a quotient graph  $R_k$  of  $M_k$ defined previously in Section 7 under action of the dihedral group  $D_{2n}$  of 2n elements.

According to Shields and Savage in their Lemma 3 [18], the existence of Hamilton paths in  $R_k$  between certain distinguished vertices yields Hamilton cycles placed in a dihedrally symmetric fashion in  $M_k$ . Thus, these paths yield the dihedrally symmetric version of Hável's conjecture mentioned above. Moreover, these paths can be presented in terms of the Kierstead-Trotter lexical matchings of  $M_k$  [13], which highlights the relevance of the said system, particularly in pursuing the existence, enumeration (more specifically, if existence is established, perhaps less than  $2^{2^{\Omega(n)}}$  paths, see next paragraph) and sorting (according to both the  $\alpha$  ordering and the lexical matchings, see Section 11) of such paths.

These tasks remain in spite of the establishment of Hável's conjecture by Mütze, who reviewed its history in [16], to which bibliography we refer, and proved that there are  $2^{2^{\Omega(n)}}$  such cycles.

More specifically, in our treatment below lexical 1-factorizations in the  $M_k$  s are defined via the Kierstead-Trotter lexical matchings [13] from Section 8 on. Ammerlaan and Vassilev [1] showed that any Hamilton cycle in a graph  $M_k$  has the same number of edges along every coordinate direction of the *n*-cube  $H_n$ . This is the case of the Shields-Savage construction of a Hamilton cycle  $\eta_k$  placed in a dihedrally symmetric fashion in  $M_k$  for it uses the cyclic nature mod *n* of  $M_k$  in taking a Hamilton path  $\xi_k$  as required and with: (i) its vertices in 1-1 correspondence with the first  $C_k$  strings  $\alpha$  and (ii) dihedral unfolding in  $M_k$  formed by concatenating  $\xi_k$  with its translates mod *n* in order to compose  $\eta_k$ , (Section 11). In fact, such  $\xi_k$  must visit just once each vertex class of  $M_k$  under an adequate equivalence relation (Section 5). The Shields-Savage construction technique (further employed in [19, 20] with no visible association to any system of numeration, naturally related, or not, to the Catalan numbers) motivates our inquiry about existence, enumeration and sorting of such Hamilton cycles  $\eta_k$ . Each vertex of  $\xi_k$  corresponds uniquely to a *k-word*  $\alpha$  in the system of numeration <u>A239903</u>, concept made precise in Section 2.

Thus, we ask: Are there as many Hamilton cycles placed in a dihedrally symmetric fashion in the  $M_k$ s as Mütze showed by his constructions in the general case in [16]? Moreover, can Mütze's construction be somehow particularized to Hamilton cycles placed in a dihedrally symmetric fashion in the  $M_k$ s?

In continuation to this approach, a generalized lexical tree is introduced in [8] whose nodes are all the k-words  $(1 \le k \in \mathbb{Z})$  and is inductively constructed in terms of corresponding k-co-words. Moreover, still in [8], the lexically colored adjacency tables in Section 10 below are analyzed in terms of k-words to contribute in answering the questions above.

### 2 Catalan-number system of numeration

According to the proposed system of numeration, the increasing sequence of non-negative integers is represented via the RGSs, starting with:

$$0, 1, 10, 11, 12, 100, 101, 110, 111, 112, 120, 121, 122, 123, 1000, \dots$$
(1)

where the subsequence 1, 10, 100, 1000, ...,  $10 \cdots 0$ , ..., (the last term shown with t zeros represented symbolically by  $10^t$ ,  $t \ge 0$ ) corresponds to the numbers  $C_1 = 1$ ,  $C_2 = 2$ ,  $C_3 = 5$ ,  $C_4 = 14, \ldots, C_{t+1} = \ldots$ , etc. Because of this, we refer to the system as the *Catalan-number* system of numeration. It can be defined formally as follows. The RGS s, representing the consecutive integers from 0 to 14, can in general be written as expressions  $a_{k-1}a_{k-2}\cdots a_2a_1$  by prefixing enough 0 s in (1) if necessary, for any adequate k. Such expressions  $a_{k-1}a_{k-2}\cdots a_2a_1$ , called k-words, are defined for  $1 < k \in \mathbb{Z}$  by means of the following two rules:

- 1. The leftmost position in a k-word  $a_{k-1}a_{k-2}\cdots a_2a_1$ , namely position k-1, contains a digit  $a_{k-1} \in \{0,1\}$ .
- 2. Given a position i > 1 with i < k in a k-word  $a_{k-1}a_{k-2}\cdots a_2a_1$ , then to the immediate right of the corresponding digit  $a_i$ , the digit  $a_{i-1}$  (meaning at position i-1) satisfies  $0 \le a_{i-1} \le a_i + 1$ .

The reader may compare this with the essentially similar *Catalan* RGSs in Section 15.2 of [2], and with the mixed radix systems [5], including the factorial number, or factoradic, system [9], [10], [14] pg. 192, [15] pg. 12, or <u>A007623</u> in [21]. We refer as well to Stanley's interpretation (u) of Catalan numbers [22], Exercise (u), as mentioned in <u>A239903</u> of [21].

Every k-word  $a_{k-1}a_{k-2}\cdots a_2a_1$  yields a (k+1)-word  $a_ka_{k-1}a_{k-2}\cdots a_2a_1 = 0a_{k-1}a_{k-2}\cdots a_2a_1$ . A k-word  $\neq 0$  stripped of the null digits to the left of the leftmost position containing digit 1 is called a *Catalan string*. We also consider (improper) *Catalan string* 0 corresponding to the null k-words,  $0 < k \in \mathbb{Z}$ .

The k-words are ordered as follows: Given any two k-words, say

$$\alpha = a_{k-1} \cdots a_2 a_1$$
 and  $\beta = b_{k-1} \cdots b_2 b_1$ ,

where  $\alpha \neq \beta$ , we say that  $\alpha$  precedes  $\beta$ , written  $\alpha < \beta$ , whenever either

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- (i)  $a_{k-1} < b_{k-1}$  or
- (ii)  $a_i = b_i$ , for  $k 1 \le j \le i + 1$ , and  $a_i < b_i$ , for some  $1 \le i < k 1$ .

The order defined on the k-words this way is said to be their *stair-wise* order.

**Observation 1.** The sequence of nonzero Catalan strings has the terms corresponding to the Catalan numbers  $C_1 = 1$ ,  $C_2 = 2$ ,  $C_3 = 5$ ,  $C_4 = 14$ , ...,  $C_{t+1} = \frac{1}{2t+3} \binom{2t+3}{t+1}$ , ..., written respectively as 1, 10, 100, 1000, ...,  $10^t$ , ..., where  $0 \le t \in \mathbb{Z}$ . Moreover, there exists exactly  $C_{k+1}$  k-words  $< 10^k$ , for each k > 0.

To determine the Catalan string corresponding to a given decimal integer  $x_0$ , or vice versa, one employs *Catalan's triangle*  $\Delta$ , namely a triangular arrangement composed by positive integers starting with the following rows  $\Delta_j$ , for  $j = 0, \ldots, 8$ :

and with a linear reading of the successive rows conforming the sequence <u>A009766</u> in [21]. Specifically, the numbers  $\tau_i^j$  in row  $\Delta_j$  of  $\Delta$  ( $0 \leq j \in \mathbb{Z}$ ) satisfy the following items:

1. 
$$\tau_0^j = 1$$
, for every  $j \ge 0$ ;  
2.  $\tau_1^j = j$  and  $\tau_j^j = \tau_{j-1}^j$ , for every  $j \ge 1$ ;  
3.  $\tau_i^j = \tau_i^{j-1} + \tau_{i-1}^j$ , for every  $j \ge 2$  and  $i = 1, \dots, j-2$ ;  
4.  $\sum_{i=0}^j \tau_i^j = \tau_j^{j+1} = \tau_{j+1}^{j+1} = C_j$ , for every  $j \ge 1$ .

A unified formula for the numbers  $\tau_i^k$ , (j = 0, 1..., k), is given by:

$$\tau_j^k = \frac{(k+j)!(k-j+1)}{j!(k+1)!}$$

Now, the determination of the Catalan string corresponding to a decimal integer  $x_0$ proceeds as follows. Let  $y_0 = \tau_k^{k+1}$  be the largest member of the second diagonal of  $\Delta$ with  $y_0 \leq x_0$ . Let  $x_1 = x_0 - y_0$ . If  $x_1 > 0$ , then let  $Y_1 = \{\tau_{k-1}^j\}_{j=k}^{k+b_1}$  be the largest set of successive terms in the (k-1)-column of  $\Delta$  with  $y_1 = \sum (Y_1) \leq x_1$ . Either  $Y_1 = \emptyset$ , in which case we take  $b_1 = -1$ , or not, in which case  $b_1 = |Y_1| - 1$ . Let  $x_2 = x_1 - y_1$ . If  $x_2 > 0$ , then let  $Y_2 = \{\tau_{k-2}^j\}_{j=k}^{k+b_2}$  be the largest set of successive terms in the (k-2)-column of  $\Delta$ with  $y_2 = \sum (Y_2) \leq x_2$ . Either  $Y_2 = \emptyset$ , in which case we take  $b_2 = -1$ , or not, in which case  $b_2 = |Y_3| - 1$ . Proceeding this way, we arrive at a null  $x_k$ . Then the Catalan string corresponding to  $x_0$  is  $a_{k-1}a_{k-2}\cdots a_1$ , where  $a_{k-1} = 1$ ,  $a_{k-2} = 1 + b_1$ , ..., and  $a_1 = 1 + b_k$ .

For example, if  $x_0 = 38$ , then  $y_0 = \tau_3^4 = 14$ ,  $x_1 = x_0 - y_0 = 38 - 14 = 24$ ,  $y_1 = \tau_2^3 + \tau_2^4 = 5 + 9 = 14$ ,  $x_2 = x_1 - y_1 = 24 - 14 = 10$ ,  $y_2 = \tau_1^2 + \tau_1^3 + \tau_1^4 = 2 + 3 + 4 = 9$ ,  $x_3 = x_2 - y_2 = 10 - 9 = 1$ ,  $y_3 = \tau_0^1 = 1$  and  $x_4 = x_3 - y_3 = 1 - 1 = 0$ , so that  $b_1 = 1$ ,  $b_2 = 2$ , and  $b_3 = 0$ , taking to  $a_4 = 1$ ,  $a_3 = 1 + b_1 = 2$ ,  $a_2 = 1 + b_2 = 3$  and  $a_1 = 1 + b_3 = 1$ , determining the 5-word of 38 to be  $a_4a_3a_2a_1 = 1231$ . If  $x_0 = 20$ , then  $y_0 = \tau_3^4 = 14$ ,  $x_1 = x_0 - y_0 = 20 - 14 = 6$ ,  $y_1 = \tau_2^3 = 5$ ,  $x_2 = x_1 - y_1 = 1$ ,  $y_2 = 0$  is an empty sum (since its possible summand  $\tau_1^2 > 1 = x_2$ ),  $x_3 = x_2 - y_2 = 1$ ,  $y_3 = \tau_0^1 = 1$  and  $x_4 = x_3 - x_3 = 1 - 1 = 0$ , determining the 5-word of 20 to be  $a_4a_3a_2a_1 = 1101$ . Moreover, if  $x_0 = 19$ , then  $y_0 = \tau_3^4 = 14$ ,  $x_1 = x_0 - y_0 = 19 - 14 = 5$ ,  $y_1 = \tau_2^3 = 5$ ,  $x_2 = x_1 - y_1 = 5 - 5 = 0$ , determining the 5-word  $a_4a_3a_2a_1 = 1100$ .

Given a Catalan string or k-word  $a_{k-1} \cdots a_1$ , the considerations above can easily be played backwards to recover the corresponding decimal integer  $x_0$ .

# **3** Descending castling

**Theorem 2.** To each k-word  $\alpha = a_{k-1} \cdots a_1$  corresponds an n-tuple  $F(\alpha)$  whose entries are the integers  $0, 1, \ldots, k$  (once each) together with k asterisks \* and such that: (a) the leftmost entry of  $F(\alpha)$  is k; (b) each integer entry to the immediate right of an integer entry b is an integer less than b; in particular,  $F(0^{k-1}) = k(k-1)(k-2) \cdots 210 * \cdots *$ ; (c) first, to each k-word  $\alpha \neq 00 \cdots 0$  corresponds a k-word  $\beta = b_{k-1} \cdots b_1$  smaller than  $\alpha$  in the stair-wise order of k-words and differing from  $\alpha$  in exactly one entry, that is  $b_i \neq a_i$  for just one i with  $k-1 \ge i \ge 1$ , being  $\beta$  maximal under these restrictions; second,  $F(\alpha) = f_0 f_1 \cdots f_{2k}$  is obtained from  $F(\beta) = g_0 g_1 \cdots g_{2k}$  by the descending castling operation consisting of:

- 1. setting  $g_0 = f_0$ ,  $g_1 = f_1$ , ...,  $g_{i-1} = f_{i-1}$ ;  $g_{2k} = f_{2k}$ ,  $g_{2k-1} = f_{2k-1}$ , ...,  $g_{2k-i+1} = f_{2k-i+1}$ ;
- 2. denoting the *i*-substrings formed by the entries in item 1 by  $W^i$  on the left and  $Z^i$  on the right, and writing  $F(\beta) \setminus (W^i \cup Z^i) = X|Y$ , (X concatenated with Y), with the substring Y starting at entry  $\ell - 1$ , where  $\ell > 0$  is the leftmost entry of X;
- 3. by noticing that  $F(\beta) = W^i |X| Y |Z^i$ , finally setting  $F(\alpha) = W^i |Y| X |Z^i$ .

In order to proceed with the proof of Theorem 2, let us note that there is a rooted tree  $\mathcal{T}_k$  whose nodes are the k-words. In fact,  $\mathcal{T}_k$  is associated to the stair-wise order of the k-words. The root of  $\mathcal{T}_k$  is  $0^{k-1}$  and any other k-word  $\alpha \neq 0^{k-1}$  is a child in  $\mathcal{T}_k$  of a k-word  $\beta$  that differs from  $\alpha$  in just one entry  $a_i$  (0 < i < k). By representing  $\mathcal{T}_k$  with the children of every k-word  $\alpha$  enclosed between parentheses after  $\alpha$  and separating siblings with commas, we can write, say for k = 4:  $\mathcal{T}_4 = 000(001, 010(011(012)), 100(101, 110(111(121)), 120(121(122(123)))))$ .

Proof. In sub-item 2 of item (c), we have: (i)  $W^i$  is a proper subsequence of the maximal starting descending integer sequence  $\overline{W}$  such that  $\overline{W} \setminus W^i$  starts with  $\ell$ , the head of X; (ii)  $Z^i$  is solely composed by asterisks. The k-tuple  $\alpha$  contains precise instructions for the recursive operation of descending castling to work as prescribed in the statement, away from those  $W^i$  and  $Z^i$  equally long in  $F(\beta)$  and  $F(\alpha)$ : By respecting the rule that in every sequence of applications of sub-items 1-3 along descending paths in  $\mathcal{T}_k$ , unit augmentation of  $a_i$  for larger values of i, (0 < i < k), must occur first, and only then in descending order of i, thus thinning the inner sub-string X|Y after each application, the resulting process effectively preserves the stated properties, for by changing the order of the appearing sub-strings X and Y, that have their first elements being respectively  $\ell$  and  $\ell - 1$  in successive decreasing order, the descending nature of this operation is effectively guaranteed.

Let us illustrate the descending castling operation. For k = 2, 3, 4, the k-words  $\alpha$  are presented in their stair-wise order in Table I, both on the left and the right columns, and their corresponding images under F, or F-images, on the penultimate column. In each of the three listings, each with  $C_k$  rows ( $C_2 = 2, C_3 = 5$  and  $C_4 = 14$ ), the columns are filled, from the second row on, as follows: (i)  $\alpha$ , appearing in downward stair-wise order; (ii)  $\beta$ as in item (c) of Theorem 2, which allows to determine the subindex i in item (iv) below, where  $\alpha$  and  $\beta$  differ; (iii)  $F(\beta)$ , from the penultimate column in the previous row; (iv) the only subindex i ( $k - 1 \ge i \ge 1$ ) for which the *i*-th entries in  $\alpha$  and  $\beta$  differ, with  $\beta$  as large as possible such that  $\beta < \alpha$ ; (v) the decomposition  $W^i|Y|X|Z^i$  of  $F(\beta)$ ; (vi) the result of the descending castling operation; (vii) the corresponding re-concatenation in the column  $F(\alpha)$ ; and (viii) again  $\alpha$ , as  $F^{-1}(F(\alpha))$ .

Clearly, for each k-word  $\alpha$  different from the zero k-word there exists a well determined  $\beta$  obtained as indicated in item (c) of Theorem 2 and observable by comparing the first two columns in Table I, which in addition highlights the sub-index *i* of the fourth column. On the other hand, not all the *n*-tuples satisfying items (a)-(b) of Theorem 2 appear in the finite

recursion implied by item (c). For example,  $\beta' = 431 * 2 * *0*$  is not image of a permissible  $\alpha$  for k = 4, (see Table I, for k = 4). In fact, trying to apply sub-items 1.-3. for i = 1 to  $\beta'$ would result in  $\alpha' = 42 * *031 * *$  which does not respect item (b). For i = 2, we would have  $\beta'$  decomposing as 43|1\*2\*\*|0\*, and since the first sub-word of the middle part starts with 1, we would like to take 0 as the second sub-word, but 0 is outside the middle part. Again this is not a case in our treatment, as indicated in item (ii) of the proof of Theorem 2.

α	$\beta$	F(eta)	i	$W^i   X   Y   Z^i$	$W^i   Y   X   Z^i$	$F(\alpha)$	α
0	_	—	—	—	—	210**	0
1	0	210**	1	2   1   0 *   *	2   0 *   1   *	20*1*	1
00	—	_	—	—	_	3210***	00
01	00	3210***	1	3   2   10**   *	3   10 * *   2   *	310**2*	01
10	00	3210***	2	32   1   0*   **	$32 \mid 0 * \mid 1 \mid * *$	320*1**	10
11	10	320*1**	1	$3 \mid 20* \mid 1* \mid *$	3   1 *   20 *   *	31*20**	11
12	11	31*20**	1	3   1*2   0*  *	$3 \mid 0 * \mid 1 * 2 \mid *$	30*1*2*	12
000	_	—	—	—	—	43210 * * * *	000
001	000	43210 * * * *	1	4 3 210 * * *  *	4 210 * * *  3 *	4210 * * * 3*	001
010	000	43210 * * * *	2	43 2 10 * *  * *	43 10 * * 2  * *	4310 * *2 * *	010
011	010	4310 * *2 * *	1	4 310 * * 2 * *	4 2* 310** *	42 * 310 * **	011
012	011	42 * 310 * **	1	4 2*3 10** *	4 10** 2*3 *	410 * *2 * 3*	012
100	000	43210 * * * *	3	432 1 0* ***	432 0* 1 ***	4320 * 1 * **	100
101	100	4320 * 1 * **	1	4 3 20*1** *	4 20*1** 3 *	420 * 1 * *3*	101
110	100	4320 * 1 * **	2	43 20* 1* **	43 1* 20* **	431 * 20 * **	110
111	110	431 * 20 * **	1	4 31 *  20 * * *	4 20** 31* *	420 * *31 * *	111
112	111	420 * *31 * *	1	4 20 * *3 1 *  *	4 1* 20**3 *	41 * 20 * *3*	112
120	110	431 * 20 * **	2	43 1*2 0* **	43 0* 1*2 **	430 * 1 * 2 * *	120
121	120	430 * 1 * 2 * *	1	4 30*1* 2* *	$4 2 \times  30 \times 1 \times   \times$	$42 \times 30 \times 1 \times 30$	121
122	121	$42 \times 30 \times 1 \times 30$	1	$4 2 \times 30 \times  1 \times   \times$	4 1 *  2 * 30 *  *	$41 \times 2 \times 30 \times 30$	122
123	122	$41 \times 2 \times 30 \times 30$	1	$4 1 \times 2 \times 3 0 \times  * $	4 0* 1*2*3 *	$40 \times 1 \times 2 \times 3 \times$	123

#### TABLE I

One may perform sub-items 1-3 departing from a k-word  $\beta$  independently of a specific  $\alpha$  by taking a subindex i in the text of item (D) and after replacing numbers and asterisks respectively by 0s and 1s, obtaining a k-work  $\alpha'$  provided already by the textual item (D). Say we depart for k = 3 from  $F(\beta) = F(01) = 310^{**}2^{*}$  and take i = 1. Sub-items 1-3 leads here to  $F(\alpha') = 30 * *21*$ , which yields 0011001, a translation mod 7 of  $31*20^{**} = F(11)$ , obtained already in a different way in Table I.

To each  $F(\alpha)$  corresponds a binary *n*-tuple  $\phi(\alpha)$  of weight k obtained by replacing each integer entry in  $\{0, 1, \ldots, k\}$  by 0 and each asterisk \* by 1. By attaching the entries of  $F(\alpha)$  as subscripts to the corresponding entries of  $\phi(\alpha)$ , a subscripted binary n-tuple  $\phi(\alpha)$ is obtained. Let  $\aleph(\phi(\alpha))$  be given by the reverse complement of  $\phi(\alpha)$ , that is

if 
$$\phi(\alpha) = a_0 a_1 \cdots a_{2k}$$
, then  $\aleph(\phi(\alpha)) = \bar{a}_{2k} \cdots \bar{a}_1 \bar{a}_0$ , (2)

where  $\overline{0} = 1$  and  $\overline{1} = 0$ . A subscripted version  $\aleph$  of  $\aleph$  is immediately obtained for  $\overline{\phi}(\alpha)$ . Observe that every image under  $\aleph$  is an *n*-tuple of weight k + 1 and has the 1s with integer subscripts and the 0s with asterisk subscripts. The integer subscripts reappear from Section 8 to Section 10 as *lexical colors* [13]. Table II illustrates the notions just presented, for k = 2, 3.

ΤA	BL	Æ	Π
TTT	DL		т.

$\alpha$	$\phi(lpha)$	$ar{\phi}(lpha)$	$\bar{\aleph}(\phi(\alpha)) = \aleph(\bar{\phi}(\alpha))$	$\aleph(\phi(\alpha))$
0	00011	$0_2 0_1 0_0 1_* 1_*$	$0_*0_*1_01_11_2$	00111
1	00101	$0_2 0_0 1_* 0_1 1_*$	$0_*1_10_*1_01_2$	01011
00	0000111	$0_3 0_2 0_1 0_0 1_* 1_* 1_*$	$0_*0_*0_*1_01_11_21_3$	0001111
01	0001101	$0_3 0_1 0_0 1_* 1_* 0_2 1_*$	$0_*1_20_*0_*1_01_11_3$	0100111
10	0001011	$0_3 0_2 0_0 1_* 0_1 1_* 1_*$	$0_*0_*1_10_*1_01_21_3$	0010111
11	0010011	$0_3 0_1 1_* 0_2 0_0 1_* 1_*$	$0_*0_*1_01_20_*1_11_3$	0011011
12	0010101	$0_3 0_0 1_* 0_1 1_* 0_2 1_*$	$0_*1_20_*1_10_*1_01_3$	0101011

An interpretation of this related to the middle-levels graphs is started at the end of Section 5 in relation to the subscripts  $0, 1, \ldots, k$  and concluded as Corollary 6 at the end of Section 9.

### 4 The middle-levels graphs

Let  $1 < n \in \mathbb{Z}$ . The *n*-cube graph  $H_n$  is the Hasse diagram of the Boolean lattice on the coordinate set  $[n] = \{0, \ldots, n-1\}$ . Vertices of  $H_n$  are cited in three different ways:

- (a) as the subsets  $A = \{a_0, a_1, \ldots, a_{r-1}\} = a_0 a_1 \cdots a_{r-1}$  of [n] they stand for, where  $0 < r \leq n$ ;
- (b) as the characteristic *n*-vectors  $B_A = (b_0, b_1, \ldots, b_{n-1})$  over the field  $\mathbb{F}_2$  that the subsets A represent, meaning they are given by  $b_i = 1$  if and only if  $i \in A$  ( $i \in [n]$ ), and represented for short by  $B_A = b_0 b_1 \cdots b_{n-1}$ ;
- (c) as the polynomials  $\epsilon_A(x) = b_0 + b_1 x + \dots + b_{n-1} x^{n-1}$  associated with the vectors  $B_A$ .

A subset A as in (a) is said to be the support of the vector  $B_A$  in (b). For each  $j \in [n]$ , the *j-level*  $L_j$  is the vertex subset in  $H_n$  formed by those  $A \subseteq [n]$  with |A| = j. For  $1 \leq k \in \mathbb{Z}$ , the *middle-levels graph*  $M_k$  is defined as the subgraph of  $H_n$  induced by  $L_k \cup L_{k+1}$ .

#### 5 Quotient graphs under cyclic action

By viewing the vertices of  $M_k$  as polynomials, as in item (c) of Section 4, an equivalence relation  $\pi$  is seen to exists in the vertex set  $V(M_k)$  of  $M_k$  by means of the logical expression:

$$\epsilon_A(x)\pi\epsilon_{A'}(x) \iff \exists i \in \mathbb{Z} \text{ such that } \epsilon_{A'}(x) \equiv x^i\epsilon_A(x) \pmod{1+x^n}.$$

Figure 1: Reflection symmetry of  $M_2/\pi$  about a line  $\ell/\pi$  and resulting graph map  $\gamma_2$ 

This implies that there exists a quotient graph  $M_k/\pi$  under a regular (i.e. transitive and free) action

$$\Upsilon': \mathbb{Z}_n \times M_k \to M_k \tag{3}$$

given by  $\Upsilon'(i, v) = v(x)x^i \mod 1 + x^n$  in polynomial terms as in item (c) of Section 4, where  $v \in V(M_k)$  and  $i \in \mathbb{Z}_n$ . Here,  $M_k/\pi$  is the graph whose vertices are the equivalence classes of vertices of  $M_k$  under  $\pi$  and whose edges are the equivalence classes that  $\pi$  induces on the edge set  $E(M_k)$  of  $M_k$ .

For example,  $M_2/\pi$  is the domain of the graph map  $\gamma_2$  suggested in Figure 1 and associated with reflective symmetry of both  $M_2/\pi$  and  $M_2$  about respective dashed vertical lines  $\ell/\pi$  and  $\ell$  acting as symmetry axes (generalized below in Section 6) with  $V(M_2/\pi) =$  $L_2/\pi \cup L_3/\pi$ , where  $L_2/\pi = \{(00011), (00101)\}$  and  $L_3/\pi = \{(00111), (01011)\}$ , each  $\pi$ -class expressed between parentheses about one of its representatives written as in (b) of Section 4 and composed by the following elements (of  $L_2/\pi$  and  $L_3/\pi$ ):

 $\begin{array}{ll} L_2/\pi = \{(00011) = \{00011, 10001, 11000, 01100, 00110\}, & (00101) = \{00101, 10010, 01001, 10100, 01010\}\}; \\ L_3/\pi = \{(00111) = \{00111, 10011, 11001, 11100, 01110\}, & (01011) = \{01011, 10101, 11010, 01101\}\}. \end{array}$ 

Let k > 1 be a fixed integer. We associate with each binary weight-k n-tuple F(A) the class (F(A)) generated by F(A) in  $L_k/\pi$  and the class  $(\aleph(F(A)))$  generated by  $\aleph(F(A))$  in  $L_{k+1}/\pi$ . These associations start the interpretation of Section 3 above to be concluded as Corollary 6 in Section 9.

#### 6 Reflective-symmetry graph involutions

A graph involution of a graph G is a graph map  $\aleph : G \to G$  such that  $\aleph^2$  is the identity graph map. Clearly, a graph involution is a graph isomorphism. In a way similar to the example for k = 2 in Section 5, but now for any  $k \geq 2$ , and in order to make explicit a graph involution  $\aleph$  of  $M_k/\pi$  given by reflective symmetry, as suggested in Figure 1 about the symmetry axis  $\ell/\pi$ , we want now to list and represent vertically the vertex parts  $L_k/\pi$  and  $L_{k+1}/\pi$  of  $M_k/\pi$  (resp.  $L_k$  and  $L_{k+1}$  of  $M_k$ ) by placing their vertices into pairs, each pair displayed on an horizontal line with its two composing vertices equidistant from a dashed vertical line  $\ell/\pi$  (resp.  $\ell$ ), like the  $\ell/\pi$  in the representation of  $M_2/\pi$  in Figure 1. To specify the desired vertex setting, the definition of  $\aleph$  in display (2) can be immediately extended to a bijection  $\aleph : L_k \to L_{k+1}$ , where the image of an element of  $L_k$  through  $\aleph$  is again said to be its reverse complement. Let us take each resulting horizontal vertex pair to be of the form  $(B_A, \aleph(B_A))$  and ordered from left to right. Let  $\rho_i : L_i \to L_i/\pi$  be the canonical projection given by assigning  $b_0 b_1 \cdots b_{n-1} \in L_i$  to  $(b_0 b_1 \cdots b_{n-1}) \in L_i/\pi$ , for i = k, k + 1, and let  $\aleph_{\pi} : L_k/\pi \to L_{k+1}/\pi$  be given by  $\aleph_{\pi}((b_0b_1\cdots b_{n-1})) = (\bar{b}_{n-1}\cdots \bar{b}_1\bar{b}_0)$ . Then  $\aleph_{\pi}$  is a bijection and we have the commutative identity  $\rho_{k+1}\aleph = \aleph_{\pi}\rho_k$ . In what follows, we say that a non-horizontal edge of  $M_k/\pi$  is a *skew edge*.

**Theorem 3.** To each skew edge  $e = (B_A)(B_{A'})$  of  $M_k/\pi$  corresponds a different skew edge  $\aleph_{\pi}((B_A))\aleph_{\pi}^{-1}((B_{A'}))$  obtained from e by reflection on the line  $\ell/\pi$ , which is equidistant from  $(B_A) = \aleph_{\pi}^{-1}((B_{A'})) \in L_k/\pi$  and  $\aleph_{\pi}((B_A)) = (B_{A'}) \in L_{k+1}/\pi$ . Thus: (i) the skew edges of  $M_k/\pi$  appear in pairs, having their endpoints in each pair forming two pairs of horizontal vertices equidistant from  $\ell/\pi$ ; (ii) the horizontal edges of  $M_k/\pi$  have multiplicity  $\leq 2$ .

Proof. With the representation adopted for the vertices of  $M_k$ , the skew edges  $B_A B_{A'}$  and  $\aleph^{-1}(B_{A'})\aleph(B_A)$  of  $M_k$  are now seen to be reflection of each other about  $\ell$  and also having their pairs  $(B_A, \aleph(B_A))$  and  $(\aleph^{-1}(B_{A'}), B_{A'})$  of endpoints lying each on a corresponding horizontal line. Now,  $\rho_k$  and  $\rho_{k+1}$  extend together to a covering graph map  $\rho : M_k \to M_k/\pi$ , since the edges accompany the projections correspondingly, as for example for k = 2, where:

$$\begin{split} &\aleph((00011)) = & \aleph(\{00011, 10001, 11000, 01100, 00110\}) = \{00111, 01110, 11100, 11001, 10011\} = (00111), \\ &\aleph^{-1}((01011)) = \aleph^{-1}(\{01011, 10110, 10110, 10101\}) = \{00101, 10010, 01001, 10100, 01010\} = (00101), \end{split}$$

showing the order of the elements in the images or preimages under  $\aleph$  of the classes mod  $\pi$  as displayed in Figure 1, that is: presented backwards (from right to left), cyclically between braces, and continuing on the right once one reaches a leftmost brace. This backwards behavior holds for any k > 2, that is:

$$\begin{split} & \otimes ((b_0 \cdots b_{2k})) = \\ & \otimes (\{b_0 \cdots b_{2k}, b_{2k} \ldots b_{2k-1}, \ldots, b_1 \cdots b_0\}) = \{\bar{b}_{2k} \cdots \bar{b}_0, \bar{b}_{2k-1} \cdots \bar{b}_{2k}, \ldots, \bar{b}_1 \cdots \bar{b}_0\} = (\bar{b}_{2k} \cdots \bar{b}_0), \\ & \otimes^{-1} ((\bar{b}'_{2k} \cdots \bar{b}'_0)) = \\ & \otimes^{-1} (\{\bar{b}'_{2k} \cdots \bar{b}'_0, \bar{b}'_{2k-1} \cdots \bar{b}'_{2k}, \ldots, \bar{b}'_1 \cdots \bar{b}'_0\}) = \{b'_0 \cdots b'_{2k}, b'_{2k-1} \cdots b'_1 \cdots b'_0\} = (b'_0 \cdots b'_{2k}), \end{split}$$

for any vertices  $(b_0 \cdots b_{2k}) \in L_k/\pi$  and  $(b'_0 \cdots b'_{2k}) \in L_{k+1}/\pi$ . This establishes item (i). Now, every edge of  $M_k$  from a vertex in  $\aleph^{-1}(v)$  to a vertex in  $\aleph^{-1}(\rho(v))$ , for some  $v \in L_k$ , projects onto an horizontal edge of  $M_k$ , while all other edges of  $M_k$  project onto corresponding skew edges of  $M_k$ . It is easy to see that an horizontal edge of  $M_k/\pi$  has its endpoint in  $L_k/\pi$ represented by a vertex  $\bar{b}_k \cdots \bar{b}_1 0 b_1 \cdots b_k \in L_k$  so there are  $2^k$  such vertices in  $L_k$  and less than  $2^k$  corresponding vertices of  $L_k/\pi$ ; for example,  $(0^{k+1}1^k)$  and  $(0(01)^k)$ , mentioned in Shields-Savage Lemma 3 [18], are endpoints of two horizontal edges, each in  $M_k/\pi$ . To prove that this implies item (ii), we have to see that there cannot be more than two representatives  $\bar{b}_k\cdots \bar{b}_1 b_0 b_1\cdots b_k$  and  $\bar{c}_k\cdots \bar{c}_1 c_0 c_1\cdots c_k$  of a vertex  $v \in L_k/\pi$ , with  $b_0 = c_0 = 0$ . Let v = $(d_0 \cdots b_0 d_{i+1} \cdots d_{j-1} c_0 \cdots d_{2k})$ , with  $b_0 = d_i$ ,  $c_0 = d_j$  and  $0 < j - i \leq k$ . A substring  $\sigma = d_{i+1} \cdots d_{j-1}$  with  $0 < j-i \le k$  is said to be (j-i)-feasible if v fulfills (ii) with multiplicity at least 2. Any (j-i)-feasible substring  $\sigma$  forces in  $L_k/\pi$  only endpoints  $\omega$  incident to two different (parallel) horizontal edges in  $M_k/\pi$  because periodic continuation mod n of  $d_0 \cdots d_{2k}$ both to the right of  $d_j = c_0$  with minimal cyclic substring  $\bar{d}_{j-1} \cdots \bar{d}_{i+1} 1 d_{i+1} \cdots d_{j-1} 0 = P_r$ and to the left of  $d_i = b_0$  with minimal cyclic substring  $0d_{i+1} \cdots d_{j-1} 1d_{j-1} \cdots d_{i+1} = P_\ell$  yields a two-way infinite string that winds up onto  $(d_0 \cdots d_{2k})$ , corresponding to  $\omega$ . For example, the initial feasible substrings  $\sigma$ , with 'o' indicating the positions  $b_0 = 0$  and  $c_0 = 0$ , are

 $(\emptyset,(oo1)), (0,(o0011)), (1,(o1o)), (0^2,(o000111)), (01,(o010011)), ((1^2,o1100)), (0^3,(o00001111)), (010,(o0100101101)), (01^2,(o0110)), (101,(o1010)), (1^3,(o111000)),$ 

where n has successive values n = 3, 5, 3, 7, 7, 5, 9, 11, 5, 5, 7. (However, the substrings  $0^{21}$  and  $10^{2}$  are non-feasible). If  $\sigma$  is a feasible substring and  $\bar{\sigma}$  is its reverse complement via  $\aleph$ ,

then the possible symmetrical substrings about  $o\sigma o = 0\sigma 0$  in (the notation of) a vertex  $\omega$  of  $L_k/\pi$  are in order of ascending length:



etc., where we use again '0' instead of 'o' for the entries immediately preceding and following the shown central copy of  $\sigma$ . Due to this, the finite lateral periods of the resulting  $P_r$  and  $P_\ell$  do not allow a third horizontal edge (at v in  $M_k$ ) up to returning to  $b_0$  or  $c_0$  since no entry  $e_0 = 0$ of  $(d_0 \cdots d_{2k})$  other than  $b_0$  or  $c_0$  happens such that  $(d_0 \cdots d_{2k})$  has a third representative  $\bar{e}_k \cdots \bar{e}_1 0 e_1 \cdots e_k$  (besides  $\bar{b}_k \cdots \bar{b}_1 0 b_1 \cdots b_k$  and  $\bar{c}_k \cdots \bar{c}_1 0 c_1 \cdots c_k$ ). Thus, those two horizontal edges are produced solely from the feasible substrings  $d_{i+1} \cdots d_{j-1}$  characterized above.  $\Box$ 

To illustrate the ideas in the proof of Theorem 3, let 1 < h < n in  $\mathbb{Z}$  be such that gcd(h, n) = 1 and let the *h*-interspersion  $\lambda_h : L_k/\pi \to L_k/\pi$  be given by  $\lambda((a_0a_1 \cdots a_n)) \to (a_0a_ha_{2h} \cdots a_{n-2h}a_{n-h})$ . For each *h* with  $1 < h \leq k$ , there is at least one *h*-feasible substring  $\sigma$  and a resulting associated vertex  $\omega \in L_k/\pi$  as in the proof of the theorem. For example, applying *h*-interspersion repeatedly by starting at  $\omega = (0^{k+1}1^k) \in L_k/\pi$  produces a number of such vertices  $\omega \in L_k/\pi$ . If we assume h = 2h' with  $h' \in \mathbb{Z}$ , then an *h*-feasible substring  $\sigma$  has the form  $\sigma = \bar{a}_1 \cdots \bar{a}_{h'}a_{h'} \cdots a_1$ , so there are at least  $2^{h'} = 2^{\frac{h}{2}}$  such *h*-feasible substrings.

### 7 Quotient graphs under dihedral action

Given a graph G with an involution  $\aleph : G \to G$ , a graph folding of G is a graph H whose vertices are the pairs  $\{v, \aleph(v)\}$ , where  $v \in V(G)$ , and whose edges are the pairs  $\{e, \aleph(e)\}$ , where  $e \in E(G)$ . Here, e has end-vertices v and  $\aleph(v)$  if and only if e yields a loop in H; otherwise, e yields a link in H ([3], pg. 3). Let us denote each pair  $((B_A), \aleph_{\pi}((B_A)))$  of  $M_k/\pi$ , horizontally represented in Section 6, via the notation  $[B_A]$ , where |A| = k.

A graph folding  $R_k$  of  $M_k/\pi$  is obtained whose vertices are the pairs  $[B_A]$  and having (1) an edge  $[B_A][B_{A'}]$  per skew-edge pair  $\{(B_A)\aleph_{\pi}((B_{A'})), (B_{A'})\aleph_{\pi}((B_A))\};$ (2) a loop at  $[B_A]$  per horizontal edge  $(B_A)\aleph_{\pi}((B_A))$ . By Theorem 3, there may be up to

(2) a loop at  $[B_A]$  per norizontal edge  $(B_A) \aleph_{\pi}((B_A))$ . By Theorem 5, there may be up to two loops at each vertex of  $R_k$ .

**Theorem 4.**  $R_k$  is a quotient graph of  $M_k$  under an action  $\Upsilon : D_{2n} \times M_k \to M_k$ .

Proof. To define  $\Upsilon$ , recall that  $D_{2n}$  is the semidirect product  $\mathbb{Z}_n \rtimes_{\varrho} \mathbb{Z}_2$  via the group homomorphism  $\varrho : \mathbb{Z}_2 \to \operatorname{Aut}(\mathbb{Z}_n)$  given by taking  $\varrho(1)$  as the automorphism assigning  $i \in \mathbb{Z}_n$  to  $(n-i) \in \mathbb{Z}_n$ , and  $\varrho((0)$  as the identity. If  $*: D_{2n} \times D_{2n} \to D_{2n}$  indicates multiplication and  $i_1, i_2 \in \mathbb{Z}_n$ , then  $(i_1, 0) * (i_2, j) = (i_1 + i_2, j)$ , but  $(i_1, 1) * (i_2, j) = (i_1 - i_2, 1 + j)$ , for  $j \in \mathbb{Z}_2$ . Now, set  $\Upsilon((i, j), v) = \Upsilon'(i, \aleph^j(v))$ , for  $i \in \mathbb{Z}_n$  and  $j \in \mathbb{Z}_2$ , where  $\Upsilon'$  was defined in display (3) of Section 5 above. It is easy to see that  $\Upsilon$  is a well-defined action of  $D_{2n}$  on  $M_k$ . For example by writing  $(i, j) \cdot v = \Upsilon((i, j), v)$  and  $v = a_0 \cdots a_{2k}$ , we have  $(i, 0) \cdot v = a_{n-i+1} \cdots a_{2k}a_0 \cdots a_{n-i} =$ v' and  $(0, 1) \cdot v' = \bar{a}_{i-1} \cdots \bar{a}_0 \bar{a}_{2k} \cdots \bar{a}_i = (n-i, 1) \cdot v = ((0, 1) * (i, 0)) \cdot v$ , leading to one instance of the compatibility condition  $((i, j) * (i', j')) \cdot v = (i, j) \cdot ((i', j') \cdot v)$  that a group action must satisfy, (together with the identity condition) to fulfill its definition.

For example, the vertices in the 20-cycle to the right in Figure 3 of Section 11 below can be rewritten as follows in their shown disposition, where a = 00011 and b = 11101:

$$\begin{split} &\Upsilon((3,0),a) \quad \Upsilon((3,1),a) \quad \Upsilon((4,0),a) \quad \Upsilon((4,1),a) \quad \Upsilon((0,0),a) \quad \Upsilon((0,1),a) \quad \Upsilon((1,0),a) \quad \Upsilon((1,1),a) \quad \Upsilon((2,0),a) \quad \Upsilon((2,1),a) \\ &\Upsilon((3,0),b) \quad \Upsilon((3,1),b) \quad \Upsilon((4,0),b) \quad \Upsilon((4,1),b) \quad \Upsilon((0,0),b) \quad \Upsilon((0,1),b) \quad \Upsilon((1,0),b) \quad \Upsilon((1,1),b) \quad \Upsilon((2,0),b) \quad \Upsilon((2,1),b) \\ &\Upsilon((2,1),b) \quad \Upsilon((2,1),b) \quad \Upsilon((2,0),b) \quad \Upsilon((2,1),b) \quad \Upsilon((2,1),b) \quad \Upsilon((2,0),b) \quad \Upsilon((2,1),b) \quad \Upsilon((2,1),b)$$

representing a Hamilton cycle in  $M_k$ , for k = 2, invariant under the action of  $D_{2n}$ .

Let the graph map  $\gamma_k : M_k/\pi \to R_k$  be the corresponding projection, as represented for k = 2 in Figure 1. Then the canonical projection  $\rho_{D_{2n}} : M_k \to R_k$  is the composition of the canonical projection  $\rho_{\mathbb{Z}_k} : M_k \to M_k/\pi$  with  $\gamma_k$ . We remark that  $\Upsilon$  is regular just when  $M_k$  is taken as a directed graph, because (undirected) edges of  $M_k$  leading to loops of  $R_k$  via  $\rho_{D_{2n}}$  appear that are fixed by  $\Upsilon$ . For example  $R_2$ , represented as the image of the graph map  $\gamma_2$  depicted in Figure 1, contains two vertices and just one (vertical) edge between them, where each vertex is incident to two loops. The representation of  $M_2/\pi$  on its left has its edges indicated with colors 0,1,2. Here, the edge  $\epsilon = (11000, 11100)$  is fixed via  $\aleph$ ; not so for the each one of the two arcs composing  $\epsilon$ .

In general, each vertex v of  $L_k/\pi$  will have its incident edges indicated with *lexical colors*  $0, 1, \ldots, k$  obtained by the following procedure arising from [13], so that  $L_k/\pi$  admits a (k+1)-edge-coloring with color set  $[k+1] = \{0, \ldots, k\}$ .

#### 8 Kierstead-Trotter lexical procedure

For each  $v \in L_k/\pi$  there are k + 1 *n*-vectors of the form  $b_0b_1 \cdots b_{n-1} = 0b_1 \cdots b_{n-1}$  that represent v with  $b_0 = 0$ . For each such *n*-vector, we take a grid  $\Gamma = P_{k+1} \Box P_{k+1}$  ([3], pg. 29), where  $P_{k+1}$  is the subgraph of the unit-distance graph of the real line  $\mathbb{R}$  induced by the set  $[k+1] \subset \mathbb{Z} \subset \mathbb{R}$ . We trace the diagonal  $\Delta$  of  $\Gamma$  from vertex (0,0) to vertex (k,k). (For k = 2,  $\Delta$  is represented by a dashed line in the instances of Figure 2, analyzed up in Section 9). Recall that an *arc* of  $\Gamma$  is an ordered pair of adjacent vertices of  $\Gamma$ . Based on [13], we build a directed 2k-path D in  $\Gamma$  from w = (0,0) to w' = (k,k) in 2k steps indexed from i = 0 to i = 2k - 1, as follows. Initially, set i = 0 and w = (0,0) and let D be the 0-path containing solely w. Repeat the following loop, formed subsequently by items (1)-(3), 2k times:

Figure 2: Representing the color assignment for k = 2

- (1) If  $b_i = 0$  (resp.,  $b_i = 1$ ), then set w' := w + (1, 0) (resp., w' := w + (0, 1)).
- (2) Augment the *i*-path D by means of the arc (w, w') of  $\Gamma$  into an (i + 1)-path, again denoted D; in other words, reset  $D := D \cup (w, w')$ ; subsequently, set i := i + 1 and w := w'.
- (3) If  $w \neq (k, k)$ , or equivalently, if i < 2k, then go back to item (1).
- (4) Set  $\bar{v} \in L_{k+1}/\pi$  as a vertex of  $M_k/\pi$  adjacent to  $v \in L_k/\pi$  and obtained from a representative  $b_0b_1\cdots b_{n-1} = 0b_1\cdots b_{n-1}$  of v by replacing the entry  $b_0$  of v by  $\bar{b}_0 = 1$  in  $\bar{v}$ , keeping the entries  $b_i$  of v with i > 0 unchanged in  $\bar{v}$ .
- (5) Set the *color* of the edge  $v\bar{v}$  to be the number c of horizontal (alternatively vertical) arcs of D below the diagonal  $\Delta$  of  $\Gamma$ .

We remark that [13] highlighted the number k + 1 - c, where c varies in [k + 1], instead of establishing a well-defined 1-1 correspondence between [k + 1] and the set of edges incident to v in  $L_k/\pi$ , as we do here. In fact, if addition and subtraction in [n] are taken modulo n and we write  $[y, x) = \{y, y + 1, y + 2, ..., x - 1\}$ , for  $x, y \in [n]$ , and  $S^c = [n] \setminus S$ , for  $S = \{i \in [n] : b_i = 1\} \subseteq [n]$ , then the cardinalities of the sets  $\{y \in S^c \setminus x : |[y, x) \cap S| < |[y, x) \cap S^c|\}$  yield all the numbers k + 1 - c in 1-1 correspondence with our colors c, where  $x \in S^c$  varies.

As in [13], the lexical procedure (or LP for short) just presented yields 1-factorizations of  $R_k$ ,  $M_k/\pi$  and  $M_k$  by means of the edge colors c = 0, 1, ..., k. This lexical approach is compatible with the graphs  $M_k/\pi$  and  $R_k$ , because each edge e of  $M_k$  has the same lexical color in [k + 1] for both arcs composing e.

### 9 Ascending castling

In what follows, a color notation  $\delta(v)$  is set for each vertex v in  $L_k/\pi$ . In fact, there exists a unique k-word  $\alpha = \alpha(v)$  with  $[F(\alpha)] = \delta(v)$ . We start by representing the lexical color assignment suggested in Figure 2 for k = 2, with the LP indicated by arrows " $\Rightarrow$ " departing from v = [00011] (top) and v = [0010] (bottom) then going to the right via depiction of working sketches of  $V(\Gamma)$  (separated by plus signs "+") for each of the three representatives  $b_0b_1\cdots b_{n-1} = 0b_1\cdots b_{n-1}$  (shown as a subtitle for each sketch, with the entry  $b_0 = 0$  underscored) in which to trace the arcs of  $D \subset \Gamma$  below  $\Delta$ , and finally pointing, via an arrow " $\rightarrow$ " departing from the representative  $b_0b_1\cdots b_{n-1} = 0b_1\cdots b_{n-1}$  in each sketch subtitle, the number of horizontal arcs of D below  $\Delta$ . Only arcs of  $D \subset \Gamma$  are traced on each sketch of  $V(\Gamma)$ , with those below  $\Delta$  indicated as darts in bold trace, and the remaining ones as segments in thin trace.

In each of the two cases in Figure 2, to the right of the three sketches, an arrow " $\Rightarrow$ " points to an unparenthesized modification of the notation  $(b_0b_1\cdots b_{n-1})$  of v obtained by setting as a subindex of each entry 0 the color obtained from its corresponding sketch, and an asterisk "\*" for each entry 1. Further to the right of this subindexed version of v, another arrow " $\Rightarrow$ " points to the string of length n formed solely by the just established subindexes in the order they appear from left to right. This final notation is indicated by  $\delta(v)$ . For each such  $\delta(v)$ , there is a unique k-word  $\alpha = \alpha(v)$  with  $(F(\alpha)) = \delta(v)$ , a fact whose proof depends on the inverse operation to descending castling in Section 3, that we may call ascending castling and that can be presented as follows:

Given an *n*-tuple v in  $L_k/\pi$ , let  $W^i = k(k-1)\cdots(k-i)$  be the maximal initial (i + 1)substring of  $\delta(v) = \delta^0(v)$ , where  $0 \le i \le k$ . If i = k, then let  $\alpha(v) = [00\cdots0]$  consist of k-1 zeros. Else, write  $\delta^0(v) = [W^i|X|Y|Z^i]$ , where  $Z^i$  is a  $j_0$ -substring with  $j_0 = i + 1$ , and X, Y start respectively at two integers  $\ell$  and  $\ell + 1 \le k - i$ . Let  $\delta^1(v) = [W^i|Y|X|Z^i]$ . If  $\delta^1(v) = [k(k-1)\cdots10*\cdots*]$ , then let  $\alpha(v)$  differ from  $a_{k-1}\cdots\alpha_1 = 00\cdots0$  just by unit incrementation of  $a_{j_0}$ . Else, repeat the procedure above starting at  $\delta^1(v)$ , and so on. In the end, we obtain a finite sequence  $\delta^0(v), \delta^1(v), \ldots, \delta^s(v)$  of *n*-tuples in  $L_k/\pi$  with parameters  $j_0 \ge j_1 \ge \ldots \ge j_s$  and k-words  $\alpha(v^0), \alpha(v^1), \ldots, \alpha(v^s) = 00\cdots0$  and obtain  $\alpha(v) = \alpha^0(v)$ from  $\alpha^s(v) = 00\cdots0 = a_{k-1}\cdots\alpha_1$  by unit incrementation of each  $a_{j_i}$ , for  $i = 0, \ldots, s$ , with each such incrementation yielding the corresponding  $\alpha(v^i)$ . Now, observe that F is a bijection between the set of k-words and the set  $L_k/\pi$ , both of cardinality  $C_k$ . This shows that in order to work with  $V(R_k)$  is enough to deal with the set of k-words, a fact useful in interpreting Theorem 5 below.

Take for example  $\delta^0(v) = [40 * 1 * 2 * 3*] = [4|0 * |1 * 2 * 3|*]$ . Then,

yielding  $\alpha(v) = \alpha(v^0) = 123$ , with the  $\alpha(v^i)$  s corresponding to the formed rows being  $\alpha(v^1) = [122], \ \alpha(v^2) = [121], \ \alpha(v^3) = [120], \ \alpha(v^4) = [110], \ \alpha(v^5) = [100] \text{ and } \ \alpha(v^6) = \alpha(v^s) = [000].$ 

Thus, the function F that sends the k-words onto their corresponding n-tuples  $F(\alpha)$  in Section 3 happens to provide a backbone relating the succeeding applications of the LP to the elements of  $L_k/\pi$ , finally covering all of  $L_k/\pi$ . A pair of skew edges  $(B_A)\aleph_{\pi}((B_{A'}))$  and  $(B_{A'})\aleph((B_A))$  in  $M_k/\pi$  is said to be a *skew reflective edge pair*. This provides a color notation for any  $v \in L_{k+1}/\pi$  such that in each particular edge class mod  $\pi$ :

- (1) each edge receives the same color regardless of the endpoint on which the LP or its modification for  $v \in L_{k+1}/\pi$  is applied;
- (2) each skew reflective edge pair in  $M_k/\pi$  is assigned a sole color in [k+1].

The modification in item (1) consists in replacing in Figure 2 each v by  $\aleph_{\pi}(v)$  so that on the left we have now instead (00111) (top) and (01011) (bottom) with respective sketch subtitles

resulting in similar sketches when the rules of the LP are taken with right-to-left readingand-processing of the entries on the left side of the subtitles (before the arrows " $\rightarrow$ "), where now the values of each  $b_i$  must be taken complemented.

Since a skew reflective edge pair in  $M_k$  determines a unique edge  $\epsilon$  of  $R_k$  (and vice versa), the color received by this pair can be attributed to  $\epsilon$ , too. Clearly, each vertex of  $M_k$  or  $M_k/\pi$  or  $R_k$  defines a bijection between its incident edges and the color set [k + 1]. The edges obtained via  $\aleph$  or  $\aleph_{\pi}$  from these edges have the same corresponding colors because of the LP.

**Theorem 5.** A 1-factorization of  $M_k/\pi$  formed by the edge colors  $0, 1, \ldots, k$  is obtained via the LP. This 1-factorization can be lifted to a covering 1-factorization of  $M_k$  and can further be collapsed onto a folding 1-factorization of  $R_k$  which induces a color notation  $\delta(v)$  on each of its vertices v. Moreover, for each  $v \in V(R_k)$  and induced notation  $\delta(v)$ , there is a unique k-word  $\alpha = \alpha(v)$  such that  $[F(\alpha)] = \delta(v)$ .

Proof. As pointed out in item (2) above, each skew reflective edge pair in  $M_k/\pi$  has its edges with the same color in [k + 1]. Thus, the [k + 1]-coloring of  $M_k/\pi$  induces a well-defined [k + 1]-coloring of  $R_k$ . This yields the claimed collapsing to a folding 1-factorization of  $R_k$ . The lifting to a covering 1-factorization in  $M_k$  is immediate. The arguments above in this section and from Section 3 determine that the collapsing 1-factorization in  $R_k$  induces the k-word  $\alpha(v)$  claimed in the statement.

**Corollary 6.**  $L_k/\pi$  and  $L_{k+1}/\pi$  can be represented respectively by the resulting classes (F(A)) and  $(\aleph(F(A)))$ .

*Proof.* The corollary follows from Theorem 5 and its preceding discussion.

#### 

## 10 Lexically colored adjacency tables

From now on and justified by Theorem 5, we use the color notation  $\delta(v)$  for the vertices v of  $R_k$  with no enclosures in parentheses or brackets as above. Furthermore, we consider a lexically colored adjacency table for  $R_k$  in which the vertices  $F(\alpha)$  of  $R_k$  are expressed via their notation  $\alpha$ , and with the order of such  $\alpha$  s taken stair-wise, as agreed before Observation 1. According to this, we view  $R_k$  as the graph whose vertices are the k-words  $\alpha$  and whose adjacency is inherited from that of their  $\delta$ -notation in  $R_k$  via pullback by  $F^{-1}$  (namely, via ascending castling). In writing elements of  $R_k$ , we avoid now parentheses or brackets.

In Table III, examples of such disposition are shown for k = 2 and 3. Notice that the neighbors of each  $F(\alpha)$  in the second column are presented as  $F^0(\alpha)$ ,  $F^1(\alpha)$ , ...,  $F^k(\alpha)$ respectively for the colors  $0, 1, \ldots, k$  of the edges incident to them, where the notation is given via the direct effect of the function  $\aleph$ . The last four columns yield the k-words  $\alpha^0, \alpha^1,$ ...,  $\alpha^k$  associated via  $F^{-1}$  respectively with the listed neighbor vertices  $F^0(\alpha)$ ,  $F^1(\alpha)$ , ...,  $F^k(\alpha)$  of  $F(\alpha)$  in  $R_k$ .

$\alpha$	$F(\alpha)$	$F^0(\alpha)$	$F^1(\alpha)$	$F^2(\alpha)$	$F^3(\alpha)$	$\alpha^0$	$\alpha^1$	$\alpha^2$	$\alpha^3$
0	210 * *	210 * *	20 * 1 *	10 * * 2	—	0	1	0	—
1	20 * 1 *	1 * 20 *	210 * *	0*1*2	_	1	0	1	—
00	3210***	3210***	320*1**	310**2*	210***3	00	10	01	00
01	310**2*	2*310**	2*30*1*	3210***	1 * 20 * * 3	01	12	00	11
10	320*1**	31*20**	3210***	30*1*2*	20*1**3	11	00	12	10
11	31*20**	320*1**	20**31*	31 *20**	10**2*3	10	11	11	01
12	30*1*2*	1*2*30*	2*310**	320*1**	0*1*2*3	12	01	10	12

TABLE III

For k = 4, observe in Table IV the resulting stair-wise adjacency disposition. In general, for any k > 1, the columns  $\alpha^i$  of the stair-wise adjacency table preserve their respective *j*-th entries (taken from right to left, so  $j = k, k - 1, \ldots, 2, 1$ ) in the following way:  $j(\alpha^0) =$  $k, j(\alpha^2) = k, j(\alpha^3) = k - 1, \ldots, j(\alpha^{k-1}) = 2, j(\alpha^k) = 1$ , while we do not have such a simple entry invariance rule for column  $\alpha^1$ . A further analysis of the relation between  $\alpha$  and each of  $\alpha^1, \ldots, \alpha^{k-1}$  proceeds in [8], Sections 2-7, in terms of k-words, as in Table IV.

#### TABLE IV

$\alpha$	$\alpha^0$	$\alpha^1$	$\alpha^2$	$\alpha^3$	$\alpha^4$
000	000	100	010	001	000
001	001	101	012	000	011
010	011	121	000	112	110
011	010	120	011	111	001
012	012	123	001	110	122
100	110	000	120	101	100
101	112	001	123	100	121
110	100	111	110	012	010
111	111	110	122	011	111
112	101	122	112	010	112
120	122	011	100	123	120
121	121	010	121	122	101
122	120	112	111	121	012
123	123	012	101	120	123
	3**	***	3**	*2*	**1

For every k > 1, each color defines an involution, as displayed in Table V, where fixed points are enclosed in parentheses, the remaining cycles are all transpositions and such fixed points and transpositions between columns  $\alpha$  and  $\alpha^i$  in the tables above are presented after a corresponding header "*i*:".

TABLE V

<i>k</i> =2	0:(0)(1); 1:(0 1); 2:(0)(1);	<i>k</i> =3	$\begin{array}{l} 0:(00)(01)(12)(10\ 11);\\ 1:(11)(00\ 10)(01\ 12);\\ 2:(11)(00\ 01)(10\ 12);\\ 3:(00)(10)(12)(01\ 11); \end{array}$
<i>k</i> =4	$\begin{array}{c} 0:(000)(001)(0\\ 1:(000)(100)(0\\ 2:(011)(110)(1\\ 3:(000)(001)(0\\ 4:(000)(100)(1\end{array}) \end{array}$	12)(111)(121)(123) 01)(101)(010 121)( 12)(121(000 010)(( 10)(112)(011 111)( 11)(112)(120)(123)	$\begin{array}{l} ((010)(011)(100\ 110)(101\ 112)(120\ 122);\\ (011\ 120)(012\ 123)(110\ 111)(112\ 122);\\ (01)(012)(100\ 120)(101\ 123)(111\ 122);\\ (012\ 110)(100\ 101)(120\ 123)(121\ 122);\\ ((001\ 011)(010\ 110)(012\ 122)(101\ 121). \end{array}$

# 11 On sorting Shields-Savage Hamilton paths

Inspired in the construction technique considered by the author and his students in [6, 7], Shields and Savage showed in their Lemma 3 [18] that a Hamilton path  $\xi_k$  in  $R_k$  starting at  $[F(0^{k-1})] = [0^{k+1}1^k]$  and ending at  $[F(12...k)] = [0(01)^k]$  exists that determines a Hamilton cycle  $\eta_k$  in  $M_k$  invariant under  $\Upsilon$  (Theorem 3 in Section 7), thus constituting a Hamilton cycle placed in a dihedrally symmetric fashion in  $M_k$ , meaning it is invariant under the dihedral-group action.

First, we pull back  $\xi_k$  via the inverse image  $\gamma_k^{-1}$  (of  $\gamma_k$  in Section 7) onto a Hamilton cycle  $\zeta_k$  in  $M_k/\pi$  invariant under  $\Upsilon'$  (Section 5), where a loop at each end of  $\xi_k$  lifts onto its corresponding parallel edge in  $M_k/\pi$ . Second, we pull  $\zeta_k$  back via  $\rho_{\mathbb{Z}_n}^{-1}$  onto a  $\eta_k$  in  $M_k$  invariant under  $\Upsilon$ .

For example, the reflection about  $\ell$  on the left of Figure 3 is used to transform  $\xi_2$  first into a Hamilton cycle  $\zeta_2$  of  $M_2/\pi$  invariant under the action of  $\mathbb{Z}_2$  induced by  $\aleph$ , represented on the figure, and then into a path of length  $2|V(R_2)| = 4$  starting at  $00101 = x^2 + x^4$  and ending at  $01010 = x + x^3$ , in the same class mod  $1 + x^5$ , that can be repeated five times to form a Hamilton cycle  $\eta_2$  invariant under  $\Upsilon$ , represented on the rest of the figure.



Figure 3: Hamilton cycles in  $M_2/\pi$  and  $M_2$ 

In the same way and again due to Lemma 3 [18], a Hamilton cycle  $\eta_k$  in  $M_k$  invariant under the action of  $D_{2n}$  is guaranteed by the determination of a Hamilton path  $\xi_k$  in  $R_k$ from vertex  $\delta(0^{k+1}1^k) = k(k-1)\cdots 21 \ast \cdots \ast$  to vertex  $\delta(0(10)^k) = k0 \ast 1 \ast 2 \ast \cdots \ast (k-1) \ast$ . A Shields-Savage Hamilton path  $\xi_k$  offers the finite sequence of colors of successive edges in  $\xi_k$  as a succinct code for the Hamilton cycle  $\eta_k$ . We can describe  $\xi_k$  via the sequence of its edge colors, or *admissible color sequence*. Table VI exemplifies in detail such sequences for k = 2, 3, namely:  $c_{1_0}^2 = 1$  and  $c_{1_0}^3 c_{2_0}^3 c_{3_0}^3 c_{4_0}^3 = 1031$ , where  $i_0$  indicates the order of presentation in which the color  $c_{i_0}^k$  is selected.

#### TABLE VI

$i_0$	$\alpha_{i_0}$	$F(\alpha_{i_0})$	$\psi(\alpha_{i_0})$	$\bar{\psi}(\alpha_{i_0})$	$\aleph(\bar{\psi}(\alpha_{i_0}))$	$\aleph(\psi(\alpha_{i_0}))$	$c_{i_0}^k$
$     \begin{array}{c}       1_{0} \\       2_{0}     \end{array} $	$\begin{array}{c} 0 \\ 1 \end{array}$	$210** \\ 20*1*$	00011 01011	$0_20_10_01_*1_* \\ 0_*1_10_*1_01_2$	$0_*0_*1_01_11_2 \\ 0_20_01_*0_11_*$	$\begin{array}{c} 00111 \\ 00101 \end{array}$	1
$10 \\ 20 \\ 30 \\ 40 \\ 50$	$\begin{array}{c} 00 \\ 10 \\ 11 \\ 01 \\ 12 \end{array}$	3210*** 320*1** 31*20** 310**2* 30*1*2*	0000111 0010111 0010011 1010011 1010010	$\begin{array}{c} 0_{3}0_{2}0_{1}0_{0}1_{*}1_{*}1_{*}\\ 0_{*}0_{*}1_{1}0_{*}1_{0}1_{2}1_{3}\\ 0_{3}0_{1}1_{*}0_{2}0_{0}1_{*}1_{*}\\ 1_{3}0_{*}1_{0}0_{*}0_{*}1_{0}1_{1}\\ 1_{*}0_{2}1_{*}0_{3}0_{0}1_{*}0_{1} \end{array}$	$\begin{array}{c} 0*0*0*1_01_11_21_3\\ 030_20_01*0_11*1*\\ 0*0*1_01_20*1_11_3\\ 0_10_01*1*0_11*0_3\\ 1_10*1_01_30*1_20* \end{array}$	$\begin{array}{c} 0001111\\ 0001011\\ 0011011\\ 0011010\\ 1011010\\ \end{array}$	$\begin{array}{c}1\\0\\3\\1\end{array}$

In fact, the feasible admissible color sequences for k = 2, 3, 4 are *lexically sorted*, that is ordered according to the Kierstead-Trotter 1-factorization, from top to bottom and then from left to right, as shown in Table VII.

#### TABLE VII

$\begin{vmatrix} \frac{k=2:}{k=3:} & \frac{1\pm}{1031+} \\ 2302- \end{vmatrix} k=4:$	$\begin{array}{c} 1013101314121\\ 1013103130121\\ 1213141310121\\ 1213143134121\\ 2140102013421\\ 3103132312031\\ 3141312312031\\ 3241041031023\\ 3241314301313\\ \end{array}$	$\begin{array}{c} 1032103132313\\ 1032141312313\\ 1214103132131\\ 1214141312131\\ 2140123423023\\ 2140142431023\\ 3230213213142\\ 3230213213142\\ 3230213231302\\ 3243130301313\\ 3243403031023\\ \end{array}$	$\begin{array}{c} 1201041413421\\ 1201314143131\\ 1214231213142\\ 121423121302\\ 2304302013421\\ 2304321021421\\ 3230303132313\\ 3230341312313\\ 3412303132131\\ 3412341312131 \end{array}$	$\begin{array}{c} 1203130143131\\ 1203403413421+\\ 1303132132413\\ 1341312132413\\ 2304323423023\\ 2304323423023\\ 2304342431023\\ 3231301310323\\ 3231303134323\\ 3431341314323-\\ 3431341314323-\\ 3431343130323\\ \end{array}$
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Indicated in the rows of Table VII with a sign  $\pm$  to the right are their admissible color sequences for which the corresponding sequences of succeeding k-words form a maximum (+) and a minimum (-) according to the stair-wise order (or  $\alpha$  ordering) of the k-words involved in the sequences. This illustrates the notions of stair-wise sorting and extremality for the admissible color sequences, to be compared with the lexical sorting and corresponding extremality in the previous paragraph. For k = 3, 4, the corresponding  $\alpha$ -ordering extrema are as follows:

 $\begin{array}{l} 00\ (1)\ 10\ (0)\ 11\ (3)\ 01\ (1)\ 12\ +\\ 00\ (2)\ 01\ (3)\ 11\ (0)\ 10\ (2)\ 12\ -\\ 000\ (1)\ 100\ (2)\ 120\ (0)\ 122\ (3)\ 121\ (4)\ 101\ (0)\ 112\ (3)\ 010\ (4)\ 110\ (1)\ 111\ (3)\ 011\ (4)\ 001\ (2)\ 012\ (1)\ 123\ +\\ 000\ (3)\ 001\ (4)\ 011\ (3)\ 111\ (1)\ 110\ (3)\ 012\ (4)\ 122\ (1)\ 112\ (3)\ 010\ (1)\ 121\ (4)\ 101\ (3)\ 100\ (2)\ 120\ (3)\ 123\ -\\ \end{array}$ 

where succeeding k-words are separated by adjacency colors expressed between parentheses, so that by concatenating from left to right the contents of those parentheses yields the corresponding admissible color sequence. On the other hand, the minimum and maximum of the admissible color sequence for k = 5 are

 $\begin{array}{l} 10102010120101043010120103421313101010121;\\ 45453545435454512545435452134242454545434. \end{array}$ 

Lexical extremality as above yields for k = 5 the contents of Table VIII.

#### TABLE VIII

0000(1) 1000(2) 1200(3) 1230(0) 1233(4) 1231(0) 1232(3)
1221 (5) 1211 (0) 1222 (1) 1120 (4) 1123 (5) 1223 (3) 1212 (4) 1010 (5) 1210 (3) 1220 (1) 1122 (4) 1121 (3) 0121 (2) 0010 (1)
$\begin{array}{c} 1010 (0) 1210 (0) 1220 (1) 1122 (1) 1121 (0) 1121 (2) 0121 (2) 0121 (1) \\ 1011 (5) 1001 (2) 1201 (1) 0112 (3) 1112 (4) 1110 (0) 1111 (3) \\ 1120 (2) 1100 (2) 1201 (4) 0101 (2) 0101 (4) 0111 (4) 0111 (3) \\ 1120 (2) 1100 (2) 1201 (4) 0101 (2) 0101 (4) 0111 (4) 0111 (3) \\ 1120 (2) 1100 (2) 1201 (4) 0101 (2) 0101 (4) 0101 (4) 0111 (4) 0111 (3) \\ 1120 (2) 1100 (2) 1201 (4) 0101 (2) 0101 (4) 0101$
0122(0)0120(3)1100(4)0101(2)0001(5)0011(4)0111(2) 0110(0)0100(4)1101(3)0123(2)0012(1)1012(2)1234 +
0000(4)0001(5)0011(0)0010(5)0110(4)0012(5)0122(2)
0112 (0) 0101 (4) 1100 (5) 0100 (4) 1101 (5) 0121 (3) 1121 (0) 1010 (3) 1000 (4) 1001 (5) 1011 (0) 1120 (4) 1123 (0) 1012 (4)
0112 (0) 0101 (4) 1100 (5) 0100 (4) 1101 (5) 0121 (3) 1121 (0) 1010 (3) 1000 (4) 1001 (5) 1011 (0) 1120 (4) 1123 (0) 1012 (4) 1210 (2) 1211 (0) 1222 (5) 1112 (0) 1122 (3) 0111 (5) 1111 (0) 1110 (2) 1221 (3) 1232 (2) 1212 (3) 1223 (4) 0120 (5) 1220 (4)

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